

Synthesis of Surveillance Strategies for Mobile Sensors

Rayna Dimitrova

University of Leicester

joint work with

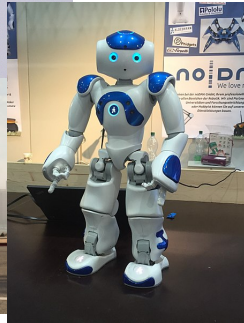
Suda Bharadwaj and Ufuk Topcu

University of Texas at Austin

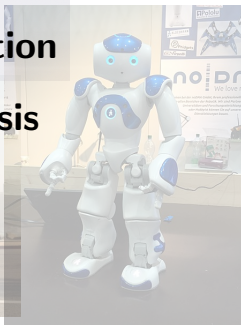
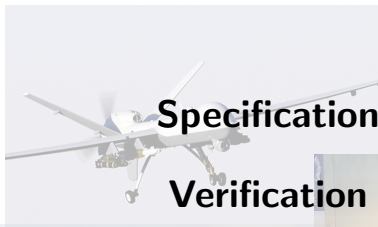
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18th September 2018

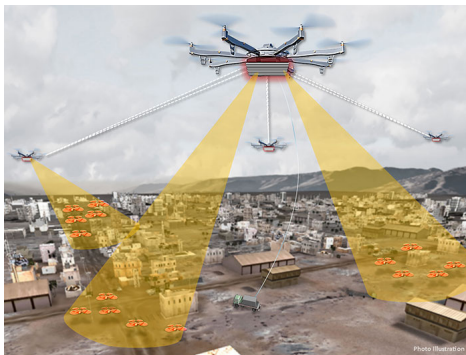
Autonomous systems: challenges and opportunities for formal methods



Autonomous systems: challenges and opportunities for formal methods



Reactive surveillance with mobile sensors



Goal: maintain knowledge of the location of a moving target

Example objectives

- ▶ always know (up to some precision) the location of the target
- ▶ eventually discover the target every time it gets out of sight

Reactive surveillance with mobile sensors

Specification φ : formulate surveillance objectives using LTL

Synthesis: solve a two player game between agent and target



agent (mobile sensor)
tries to satisfy φ



target
tries to violate φ

Compute a strategy for the agent to enforce φ .

Reactive surveillance with mobile sensors

Specification φ : formulate surveillance objectives using LTL

- ▶ introduce **surveillance predicates**

Synthesis: solve a two player game between agent and target

- ▶ tracking agent's **knowledge**
- ▶ handling **multiple sensors**

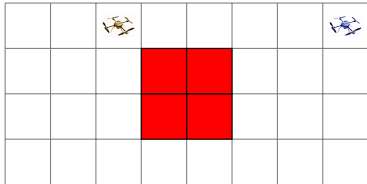
"Synthesis of Surveillance Strategies via Belief Abstraction"

S. Bharadwaj, **R. D.**, U. Topcu, CDC 2018

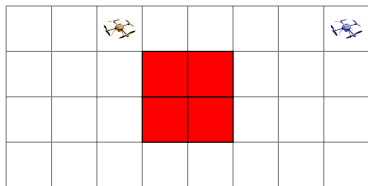
"Distributed Synthesis of Surveillance Strategies for Mobile Sensors"

S. Bharadwaj, **R. D.**, U. Topcu, CDC 2018

Surveillance game structures

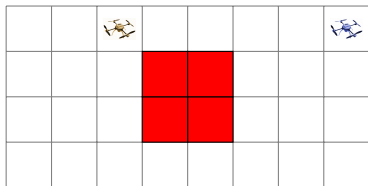


Surveillance game structures



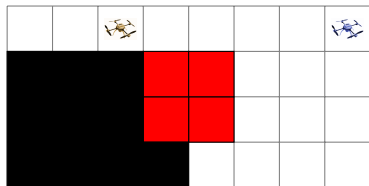
► set of locations L

Surveillance game structures



- ▶ set of locations L
- ▶ states $(l_a, l_t) \in L \times L$
 - l_a : location of agent
 - l_t : location of target

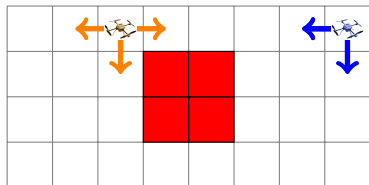
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visibility: $vis(l_a, l_t) = true$ iff l_t is in the line of sight of l_a

Surveillance game structures

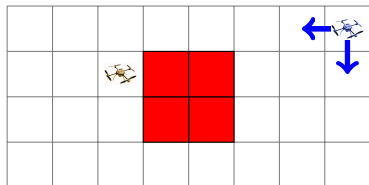


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- ▶ transitions $T, (l_a, l_t) \rightsquigarrow (l'_a, l'_t)$

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transitions: move of target, followed by move of agent

Surveillance game structures

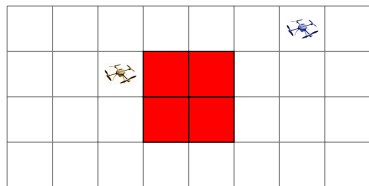


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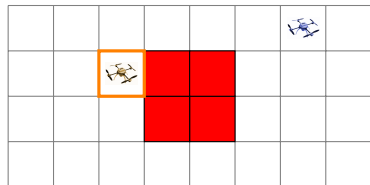
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Surveillance objectives

Belief game structure

belief: knowledge about the possible current locations of target



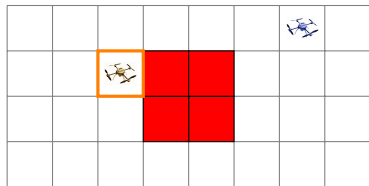
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▶ belief states $(l_a, B_t) \in L \times 2^L$

Surveillance objectives

Belief game structure

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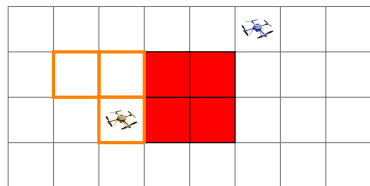
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belief transitions track the evolution of the agent's belief

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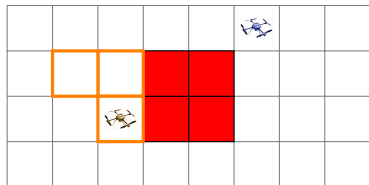
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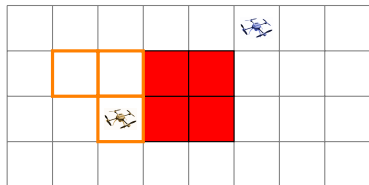
Specification

belief predicate $p_{\leq b}$, for $b \in \mathbb{N}_{>0}$: $(l_a, B_t) \models p_{\leq b}$ iff $|B_t| \leq b$

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LTL surveillance formulas: LTL with belief predicates. Examples:

- ▶ **safety surveillance** $\square p_{\leq b}$: "always" $p_{\leq b}$
- ▶ **liveness surveillance** $\square \diamond p_{\leq b}$: "infinitely often" $p_{\leq b}$

Surveillance games and strategies

surveillance game (G, φ) , where

- ▶ $G = (L, vis, T)$ is a surveillance game structure,
- ▶ φ is a surveillance specification

strategy for the agent: function that maps sequences of belief states to moves that agree with T

A strategy for the agent is **winning** in (G, φ) if each sequence of belief states resulting from this strategy satisfies the specification φ .

Synthesis of surveillance strategies

Surveillance synthesis problem

Given: surveillance game (G, φ)

Compute: strategy for the agent winning in (G, φ)

A possible approach:

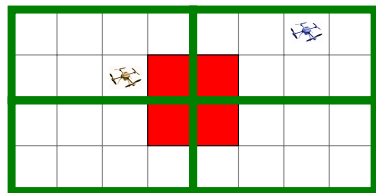
Solve game with LTL objective over belief game structure

Problem:

Size of belief game structure can be exponential in $|L|$

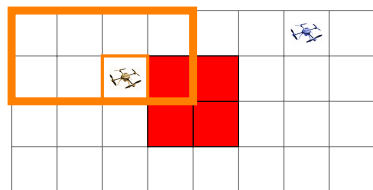
⇒ **Use abstraction!**

Belief abstraction



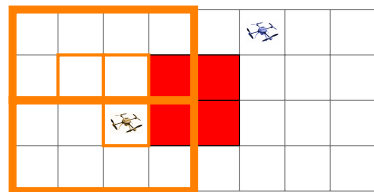
- ▶ $Q = \{Q_i\}_{i=1}^n$ partition L
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Belief abstraction



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 $(l_a, A_t) \in L \times (2^Q \cup L)$

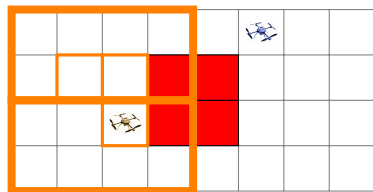
Belief abstraction



- ▶ $\mathcal{Q} = \{Q_i\}_{i=1}^n$ partition L
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abstract belief transition: overapproximate belief at each step

Belief abstraction

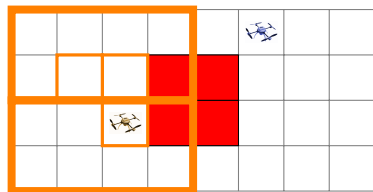


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Belief abstraction is sound for surveillance objectives.

Belief abstraction



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Belief abstraction is sound for surveillance objectives.

Worst case abstraction: each Q_i is singleton.

Abstraction-based synthesis of surveillance strategies

Abstract surveillance game: two-player game with LTL objective

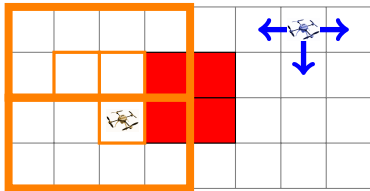
⇒ use methods for synthesis of reactive systems

Restrict surveillance objectives to the efficient fragment GR(1)

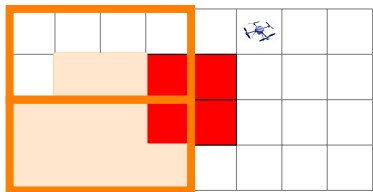
⇒ use slugs [Ehlerers and Raman 2016]

Winning abstract strategy for agent \mapsto surveillance strategy

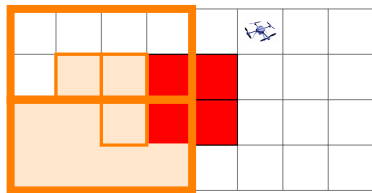
Abstract counterexamples



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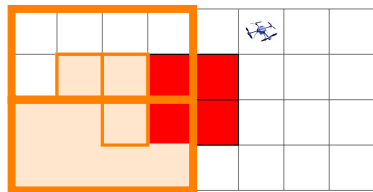


Abstract counterexamples



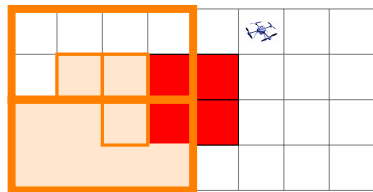
- ▶ specification $\square p_{\leq 2}$
⇒ concretizable

Abstract counterexamples



- ▶ specification $\square p_{\leq 2}$
 \Rightarrow concretizable
- ▶ specification $\square p_{\leq 5}$
 \Rightarrow spurious

Abstract counterexamples

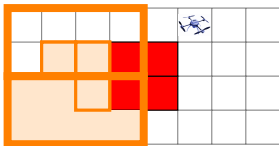


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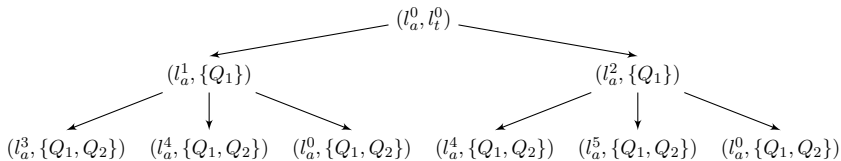
Analyse counterexample by computing concrete beliefs.

Determine which partitions to split, to refine the belief abstraction.

Counterexample-based belief refinement



abstract counterexample for the surveillance specification $\square p_{\leq 5}$

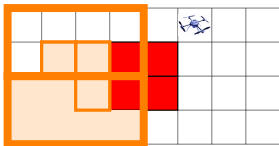


Annotate nodes of the tree with concrete belief sets.

Check if there is a leaf node where the bound is not exceeded.

If yes, then the counterexample is spurious. Refine to eliminate it.

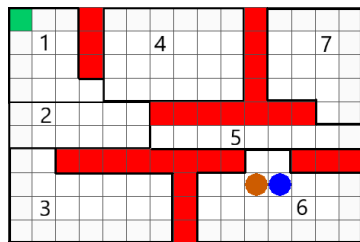
Counterexample-based belief refinement



Counterexamples for general surveillance properties are finite graphs.

- ▶ For a liveness property $\square \diamond p_{\leq b}$, check if there is a lasso path with a concrete belief in the loop with size not exceeding b .
- ▶ For general properties: refine some node with imprecise belief.

Example with liveness surveillance objective



specification

$\square \diamond p_{\leq 1} \wedge \square \diamond goal$

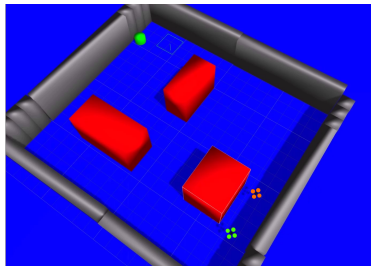
mobile sensor

straight-line visibility up to 5 cells

Number of abstract belief sets $15 \cdot 10 + 2^7$

Number of concrete belief sets 2^{150}

Example with safety surveillance objective



specification

$\square p_{\leq 30} \wedge \square \diamond goal$

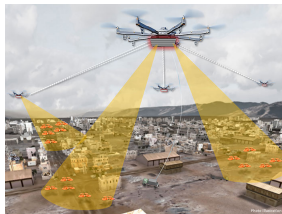
mobile sensor

unbounded straight-line visibility

Number of abstract belief sets $13 \cdot 18 + 2^6$

Number of concrete belief sets $\approx 2^{234}$

Multiple sensors



In practice: multiple sensors

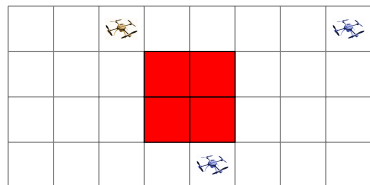


better coverage, smaller abstractions should suffice



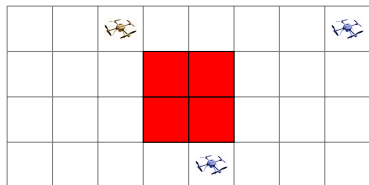
the size of the state space of the concrete game increases

Multi-agent surveillance game structures



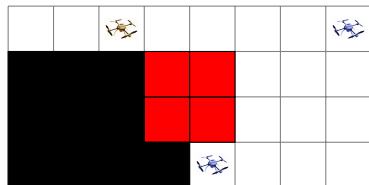
► set of locations L

Multi-agent surveillance game structures



- ▶ set of locations L
- ▶ states $(l_a^1, \dots, l_a^m, l_t)$

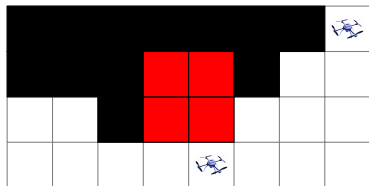
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visibility: $vis_i(l_a^i, l_t) = true$ iff l_t is in the line of sight of l_i

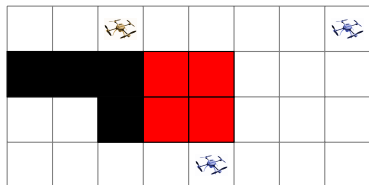
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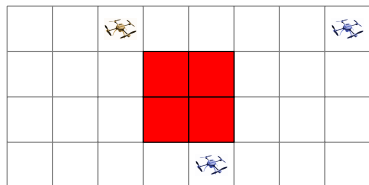


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- ▶ joint visibility $vis : L^{m+1} \rightarrow \mathbb{B}$

visibility: $vis_i(l_a^i, l_t) = true$ iff l_t is in the line of sight of l_i

joint visibility: $vis(\vec{l}, l_t) = true$ iff l_t is visible to at least one agent

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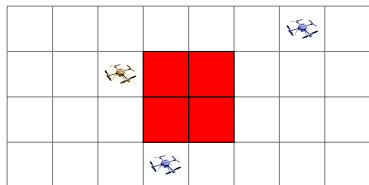
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transitions: move of target, followed by agents' synchronous move

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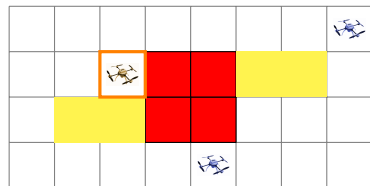
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Multi-agent surveillance with static sensors

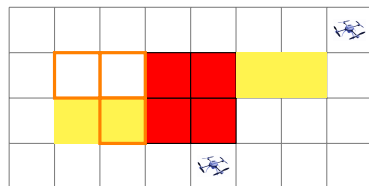


► static sensors (R_1, \dots, R_k)

static sensor: defined by its range $R_i \subseteq L$

Static sensors do not exhibit false positives or false negatives.

Multi-agent surveillance with static sensors



▶ static sensors (R_1, \dots, R_k)

▶ belief states

$$(l_a, B_t, C) \in L \times 2^Q \times 2^{\{1, \dots, k\}}$$

static sensor: defined by its range $R_i \subseteq L$

Static sensors do not exhibit false positives or false negatives.

B_t is contained in the ranges of the triggered sensors C .

Multi-agent surveillance strategies

multi-agent surveillance game $(G, \{R_1, \dots, R_k\}, \varphi)$, where

- ▶ G is a multi-agent surveillance game structure,
- ▶ R_1, \dots, R_k are static sensors,
- ▶ φ is a surveillance specification

A **joint strategy** for the agents is **winning** in $(G, \{R_1, \dots, R_k\}, \varphi)$ if each sequence of belief states resulting from the strategies for the agents satisfies the specification φ .

Multi-agent surveillance strategy synthesis

Multi-agent surveillance synthesis problem

Given: multi-agent surveillance game $(G, \{R_1, \dots, R_k\}, \varphi)$

Compute: joint strategy for the agents that is winning

A possible approach:

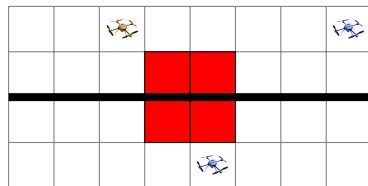
Compute a centralized strategy.

Problem:

Size of the state space is exponential in m .

⇒ **Decompose the synthesis problem!**

Game structure decomposition

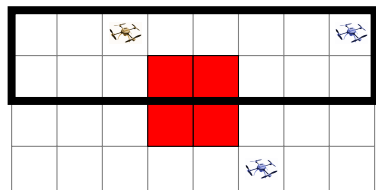


- ▶ partition $L = L_1 \uplus \dots \uplus L_m$
- ▶ agent i cannot exit L_i
- ▶ agent i cannot observe $L \setminus L_i$

Synthesize individual surveillance strategies independently.

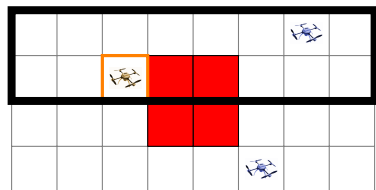
Define local specifications appropriately to ensure soundness.

Local surveillance game structures



- ▶ locations $L_i \uplus \{\hat{l}_i\}$
- ▶ states (l_a, l_t)
- ▶ $\widehat{vis}_i(l_a, \hat{l}_i) = false$
- ▶ static sensors \hat{R}_i

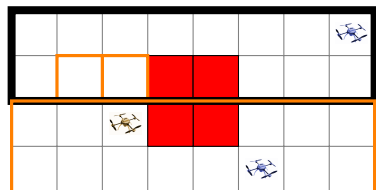
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- ▶ locations $L_i \uplus \{\hat{l}_i\}$
- ▶ states (l_a, l_t)
- ▶ $\widehat{vis}_i(l_a, \hat{l}_i) = false$
- ▶ static sensors \hat{R}_i

Agent 1: size of local belief set is 1

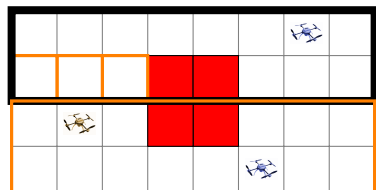
Local surveillance game structures



- ▶ locations $L_i \uplus \{\hat{l}_i\}$
- ▶ states (l_a, l_t)
- ▶ $\widehat{vis}_i(l_a, \hat{l}_i) = false$
- ▶ static sensors \hat{R}_i

Agent 1: size of local belief set is 3, including \hat{l}_1

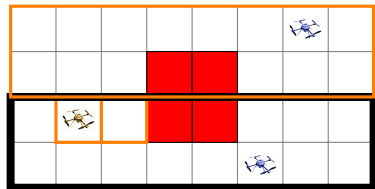
Local surveillance game structures



- ▶ locations $L_i \uplus \{\hat{l}_i\}$
- ▶ states (l_a, l_t)
- ▶ $\widehat{vis}_i(l_a, \hat{l}_i) = false$
- ▶ static sensors \hat{R}_i

Agent 1: size of local belief set is 4, including \hat{l}_1

Local surveillance game structures

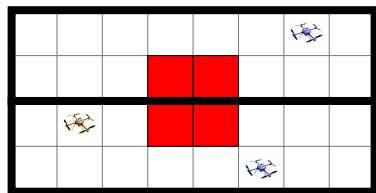


- ▶ locations $L_i \uplus \{\hat{l}_i\}$
- ▶ states (l_a, l_t)
- ▶ $\widehat{vis}_i(l_a, \hat{l}_i) = false$
- ▶ static sensors \hat{R}_i

Agent 1: size of local belief set is 4, including \hat{l}_1

Agent 2: size of local belief set is 3, including \hat{l}_2

Local surveillance game structures



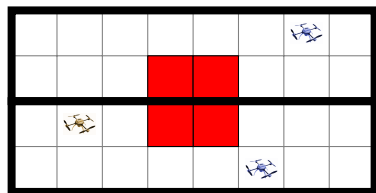
- ▶ locations $L_i \uplus \{\hat{l}_i\}$
- ▶ states (l_a, l_t)
- ▶ $\widehat{vis}_i(l_a, \hat{l}_i) = false$
- ▶ static sensors \hat{R}_i

Agent 1: size of local belief set is 4, including \hat{l}_1

Agent 2: size of local belief set is 3, including \hat{l}_2

The size of the global belief set is 5.

Global belief sets



- ▶ locations $L_i \uplus \{\widehat{l}_i\}$
- ▶ states (l_a, l_t)
- ▶ $\widehat{vis}_i(l_a, \widehat{l}_i) = false$

local belief set of agent i : $\widehat{B}_t^i \subseteq (L_i \uplus \{\widehat{l}_i\})$

global belief set of agent i : $B_t^i = \begin{cases} \widehat{B}_t^i & \text{if } \widehat{l}_i \notin \widehat{B}_t^i \\ \widehat{B}_t^i \cup (L \setminus L_i) & \text{otherwise} \end{cases}$

joint global belief set: $\bigcap_{i \in \{1, \dots, m\}} B_t^i$

Specification decomposition

We want **local surveillance specifications** $\varphi_1, \dots, \varphi_n$ such that if f_1, \dots, f_n are winning strategies in the local games $(G_i, \widehat{R}_i, \varphi_i)$ then $f_1 \otimes \dots \otimes f_n$ is a winning strategy in $(G, \{R_1, \dots, R_k\}, \varphi)$.

Specification decomposition

We restrict to conjunctions of safety and liveness surveillance.

$$\Box p_{\leq a} \wedge \Box p_{\leq b} \quad \equiv \quad \Box p_{\leq \min(a,b)}$$

$$\Box \Diamond p_{\leq a} \wedge \Box \Diamond p_{\leq b} \quad \equiv \quad \Box \Diamond p_{\leq \min(a,b)}$$

$$\Box p_{\leq a} \wedge \Box \Diamond p_{\leq b} \quad \equiv \quad \Box p_{\leq a} \quad \text{if } a \leq b$$

It suffices to consider only specifications of the following forms

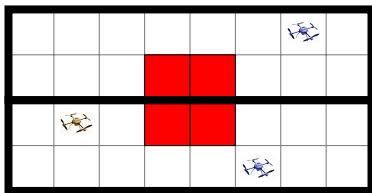
- ▶ safety $\Box p_{\leq a}$, liveness $\Box \Diamond p_{\leq a}$,
- ▶ mixed $\Box p_{\leq a} \wedge \Box \Diamond p_{\leq b}$ with $a > b$.

Safety surveillance objectives

For global specification $\square p_{\leq b}$ and $n \geq 2$ agents, take local specifications

$$\square p_{\leq c}, \text{ where } c = \lfloor \frac{b}{n} \rfloor + 1.$$

Example: specification $\square p_{\leq 2}$



Each of the local specifications is $\square p_{\leq 2}$ as well.

Conservative approximation due to the absence of coordination.

Liveness surveillance objectives

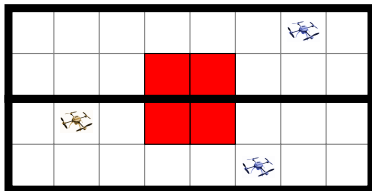
Require that each mobile sensor satisfies the liveness specification.

For global specification $\Box\Diamond p_{\leq 2}$ and n agents, take

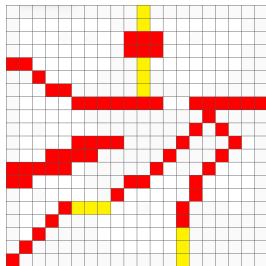
$$(\Box\Diamond(\textit{belief} \neq \{\widehat{l}_i\})) \rightarrow (\Box\Diamond(p_{\leq b} \wedge (\widehat{l}_i \notin \textit{belief}))),$$

where $\textit{belief} \neq \{\widehat{l}_i\}$ and $\widehat{l}_i \notin \textit{belief}$ are surveillance predicates.

Example: specification $\Box\Diamond p_{\leq 1}$



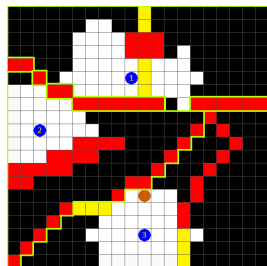
Example



- ▶ model terrain by 20×20 grid
- ▶ red regions: impassable terrain
- ▶ yellow regions: range of static sensors

Surveillance specification: $\square \diamond p_{\leq 5}$

Example



	Subgame	Number of locations	Synthesis time (s)
3 sensors	Subgame 1	142	473
	Subgame 2	113	306
	Subgame 3	145	372
	Total	400	1151
6 sensors	Subgame 1	69	101
	Subgame 2	74	206
	Subgame 3	62	111
	Subgame 4	52	88
	Subgame 5	77	285
	Subgame 6	66	64
	Total	400	855

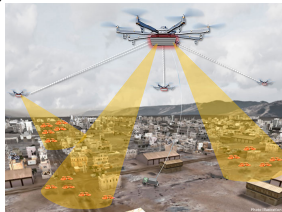
- ▶ model terrain by 20×20 grid
- ▶ red regions: impassable terrain
- ▶ yellow regions: range of static sensors

Surveillance specification: $\square \diamond p_{\leq 5}$

Current work and future directions

- ▶ Heuristics for constructing initial abstraction
- ▶ Improved abstraction refinement methods
- ▶ Less conservative specification decomposition
- ▶ Some coordination between mobile sensors
- ▶ Probabilistic detection errors by static sensors
- ▶ Noisy observations from mobile sensors

...



Conclusion

- ▶ Applying reactive synthesis to surveillance problems
- ▶ Domain specific formal specification languages
- ▶ Customized abstraction and refinement methods
- ▶ Compositional approaches key for achieving scalability

Thank you for your attention!



Questions?

Papers at 57th IEEE Conference on Decision and Control
preprints available at [raynadimitrova.github.io](https://github.com/raynadimitrova)