Graphs without large bicliques and well-quasi-orderability by the induced subgraph relation*

AISTIS ATMINAS, VADIM LOZIN, AND IGOR RAZGON

Recently, Daligault, Rao and Thomassé asked in [3] if every hereditary class which is well-quasi-ordered by the induced subgraph relation is of bounded clique-width. While the question has been shown to have a negative answer in general [9], in the present paper we show that the statement is true for a family of hereditary classes of graphs that exclude large bicliques as subgraphs. In particular, this implies (through the use of Courcelle theorem [2]) that any problem definable in Monadic Second Order Logic can be solved in a polynomial time for all well-quasi-ordered hereditary classes of graphs that exclude large bicliques.

1. Introduction

Well-quasi-ordering is a highly desirable property and a frequently discovered concept in mathematics and theoretical computer science [6, 8]. One of the most remarkable recent results in this area is the proof of Wagner's conjecture stating that the set of all finite graphs is well-quasi-ordered by the minor relation [12]. However, the subgraph or induced subgraph relation is not a well-quasi-order. On the other hand, each of these relations may become a well-quasi-order when restricted to graphs with some special properties.

A graph property (or a class of graphs) is a set of graphs closed under isomorphism. A property is hereditary if it is closed under taking induced subgraphs. It is well-known (and not difficult to see) that a graph property X is hereditary if and only if X can be described in terms of forbidden induced subgraphs. More formally, X is hereditary if and only if there is a set M of graphs such that no graph in X contains any graph from M as an induced subgraph. We call M the set of forbidden induced subgraphs for X and say that the graphs in X are M-free.

^{*}This works was supported by EPSRC grant EP/L020408/1.

Of our particular interest in this paper are graphs without large bicliques. We say that the graphs in a hereditary class X are without large bicliques if there is a natural number t such that no graph in X contains $K_{t,t}$ as a (not necessarily induced) subgraph. Equivalently, there are q and r such $K_{q,q}$ and K_r appear in the set of forbidden induced subgraphs for X. According to [11], these are precisely the graphs with a subquadratic number of edges. This family of properties includes many important classes, such as graphs of bounded vertex degree, of bounded tree-width, all proper minor closed graph classes. In all these examples, the number of edges is bounded by a linear function in the number of vertices and all of the listed properties are rather small (see e.g. [10] for the number of graphs in proper minor closed graph classes). In the terminology of [1], they all are at most factorial. In fact, the family of classes without large bicliques is much richer and contains classes with a superfactorial speed of growth, such as projective plane graphs (or more generally C_4 -free bipartite graphs), in which case the number of edges is $\Theta(n^{\frac{3}{2}})$.

Recently, Daligault, Rao and Thomassé asked in [3] if every hereditary class which is well-quasi-ordered by the induced subgraph relation is of bounded clique-width. While the question has been shown to have a negative answer in general [9], the relationship holds true for some families of hereditary graph classes. Investigating such families is interesting because it connects two seemingly unrelated notions and leads to a strong algorithmic consequence. Indeed, it follows (through the use of Courcelle theorem [2]) that for such families any problem definable in Monadic Second Order Logic can be solved in a polynomial time on any class well-quasi-ordered by the induced subgraph relation.

In the present paper, we establish the relationship between well-quasiordering and boundedness of clique-width for graphs without large bicliques. More precisely, we prove that if a class X without large bicliques is wellquasi-ordered by the induced subgraph relation, then the graphs in X have bounded path-width, i.e. there is a constant c such that the path-width of any graph in X is at most c. Since bounded path-width implies bounded clique-width, the result affirmatively answers the question in [3] for graphs without large bicliques. Thus the above algorithmic consequence is confirmed e.g. for classes of graphs of bounded degree.

Section 2 contains all preliminary information related to the topic. In this section we define an infinite family of graphs pairwise incomparable by the induced subgraph relation, which we call *canonical graphs*. In Section 3 we prove our main combinatorial result, Theorem 1, stating that a graph without large bicliques and having a large path-width has a large induced

canonical graph. A consequence of this result is that if a class X without large bicliques has unbounded path-width, then X contains an infinite subset of canonical graphs, i.e. an infinite antichain. This implies that classes of graphs without large bicliques that are well quasi-ordered by the induced subgraph relation must have bounded path-width.

2. Notation and definitions

In this work we will be using standard graph theory terminology and notation consistent with the book of Diestel [4]. In particular, K_n and P_n denote the complete graph and the chordless path with n vertices, respectively, and $K_{n,m}$ stands for a complete bipartite graph with parts of size n and m.

Throughout the text, whenever we say that G contains H, we mean that H is a subgraph of G, unless we explicitly say that H is an *induced* subgraph of G (or G contains H as an *induced* subgraph). If H is not an induced subgraph of G, we say that G is H-free. By R = R(k, r, m), we denote the Ramsey number, i.e. the minimum R such that in every colouring of k-subsets of an R-set with r colours there is a monochromatic m-set, i.e. a set of m elements all of whose k-subsets have the same colour.

According to the celebrated Graph Minor Theorem of Robertson and Seymour, the set of all graphs is well-quasi-ordered by the graph minor relation [12]. This, however, is not the case for the more restrictive relations such as subgraph or induced subgraph. Indeed, a sequence of graphs $H_1, H_2, ...$, creates an infinite antichain with respect to both relations, where H_i is the graph represented in Figure 1.



Figure 1: The graph H_i .

By connecting two vertices of degree one having a common neighbour in H_i , we obtain a graph represented on the left of Figure 2. Let us denote this graph by H'_i . By further connecting the other pair of vertices of degree one we obtain the graph H''_i represented on the right of Figure 2.

We call any graph of the form H_i , H'_i or H''_i an H-graph. Furthermore, we will refer to H''_i a $tight\ H$ -graph and to H'_i a $semi-tight\ H$ -graph. In an

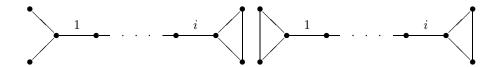


Figure 2: Graphs H'_i and H''_i .

H-graph, the path connecting two vertices of degree 3 will be called the *body* of the graph, and the vertices which are not in the body the *wings*.

Following standard graph theory terminology, we call a chordless cycle of length at least four a *hole*. Let us denote by

 \mathcal{C} the set of all holes and all H-graphs.

It is not difficult to see that any two distinct (i.e. non-isomorphic) graphs in \mathcal{C} are incomparable with respect to the induced subgraph relation. In other words,

Claim 1. C is an antichain with respect to the induced subgraph relation.

Moreover, from the poof of Theorem 1 we will see that for classes of graphs without large bicliques which are of unbounded path-width this antichain is unavoidable, or *canonical*, in the terminology of [5]. Suggested by this observation, we introduce the following definition.

Definition 1. The graphs in the set C will be called CANONICAL.

The *order* of a canonical graph G is either the number of its vertices, if G is a hole, or the the number of vertices in its body, if G is an H-graph.

3. Main result

In this section we prove the following theorem which is the main result of the paper.

Theorem 1. If X is a hereditary subclass of $(K_t, K_{q,q})$ -free graphs which is well-quasi-ordered by the induced subgraph relation, then graphs in X have a bounded path-width.

To prove the theorem, we will show that a large path-width combined with the absence of large bicliques implies the existence of a large induced canonical graph, which is a much richer structural consequence than just the existence of a long induced path. An important part of showing the existence of a large canonical graph is verifying that its body (see Section 2 for the terminology) is induced. This will be done by application of the following theorem proved in [7].

Theorem 2. For every s, t, and q, there is a number Z = Z(s,t,q) such that every graph with a path of length at least Z contains either P_s or K_t or $K_{q,q}$ as an induced subgraph.

A plan of the proof of Theorem 1 is outlined in Section 3.1. Sections 3.2, 3.3, 3.4, 3.5 contain various parts of the proof.

3.1. Plan of the proof

To prove Theorem 1 we will show that graphs of arbitrarily large pathwidth contain either arbitrarily large bicliques as subgraphs or arbitrarily large canonical graphs as induced subgraphs. The main notion in our proof is that of a *rake-graph*.

A rake-graph (or simply a rake) consists of a chordless path, the base of the rake, and a number of pendant vertices, called teeth, each having a private neighbour on the base. The only neighbour of a tooth on the base will be called the root of the tooth, and a rake with k teeth will be called a k-rake. We will say that a rake is ℓ -dense if any ℓ consecutive vertices of the base contain at least one root vertex. An example of a 1-dense 9-rake is given in Figure 3.

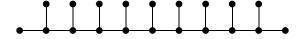


Figure 3: 1-dense 9-rake.

We will prove Theorem 1 through a number of intermediate steps as follows.

- 1. In Section 3.2, we observe that any graph of large path-width contains a rake with many teeth as a subgraph.
- 2. In Section 3.3, we show that any graph containing a rake with many teeth as a subgraph contains either
 - a dense rake with many teeth as a subgraph or
 - a large canonical graph as an *induced* subgraph.

- 3. In Section 3.4, we prove that dense rake subgraphs necessarily imply either
 - a large canonical graph as an *induced* subgraph or
 - a large biclique as a subgraph.
- 4. In Section 3.5, we use the results of sections 3.2–3.4 to deduce Theorem 1.

3.2. Rake subgraphs in graphs of large path-width

Lemma 1. For any natural k, there is a number f(k) such that every graph of path-width at least f(k) contains a k-rake as a subgraph.

Proof. In [13], Robertson and Seymour has shown that for any tree T there is a constant c_T such that any graph of path-width is at least c_T contains T as a minor. Taking T to be some fixed k-rake, we obtain that there exist a constant f(k) such that any graph of path-width at most f(k) contains a k-rake as a minor. Finally, it is not hard to see that if a graph contains a k-rake as a minor, then it also contains a k-rake as a subgraph. This observation completes the proof.

3.3. From rake subgraphs to dense rake subgraphs

Lemma 2. Let k and s be natural numbers. Every graph containing a k+2-rake as a subgraph contains either

- an s + 5-dense k-rake as a subgraph or
- a canonical graph of order at least s as an induced subgraph.

Proof. Consider a graph that contains a k+2-rake as a subgraph and choose such a k+2-rake with the minimal number of vertices. We denote the base of the rake by P. Let $\{u_1, u_2, \ldots, u_{k+2}\}$ denote the roots of the rake that are indexed respecting the linear order of the path P, i.e. so that u_1 and u_{k+2} are the endpoints of P and the subpaths of P from u_i to u_{i+1} , which we denote by P_i , are all mutually disjoint apart from the endpoints. Note that by minimality of the rake it follows that each endpoint of the path P is indeed a root vertex of the rake and that each P_i is an induced path. If each P_i for $i=2,3,\ldots,k$ has at most s+5 vertices, then we have an s+5-dense k-rake as required. So assume now that P_i for some $i=2,3,\ldots k$ has size more than s+5. To complete the proof we will show that this P_i gives rise to a canonical graph of order at least s as an induced subgraph. We proceed with some notation.

Let $P_i = w_1 w_2 \dots w_r$ with $w_1 = u_i$ and $w_r = u_{i+1}$. Extend P_i by adding the vertex w_0 of P_{i-1} that is adjacent to w_1 and the vertex w_{r+1} of P_{i+1} that is adjacent to w_r (unique choice as P_{i-1} and P_{i+1} are induced paths). Note that $w_0 w_1 w_2 \dots, w_r w_{r+1}$ is a subpath of P, the tooth v_i is adjacent to w_1 and the tooth v_{i+1} is adjacent to w_r . Let G be a graph induced by vertices $\{w_0, w_1, \dots, w_{r+1}\} \cup \{v_i, v_{i+1}\}$ and note that G contains an H-graph formed by edges $\{w_0 w_1, w_1 w_2, \dots, w_r w_{r+1}\} \cup \{v_i w_1, v_{i+1} w_r\}$ as a subgraph but not necessarily as an induced subgraph. Note that the body of the H-graph, spanned by vertices $\{w_1, w_2 \dots, w_r\}$, is a chordless path P_i . For the rest of the proof we will be arguing on the adjacencies of the wings of the H-graph in G, i.e. adjacencies of vertices w_0, w_r, v_i and v_{i+1} in G. It will follow G contains a canonical subgraph of order at least s as an induced subgraph.

We first claim that w_0 is not adjacent to w_l for any $l=2,3,\ldots,r-1$. Indeed, suppose for contradiction that w_0 is adjacent to some w_l for $l=2,3,\ldots,r-1$. Let a path P' be obtained from path P by replacing subpath $w_0w_1\ldots w_r$ of P by path $w_0w_lw_{l+1}\ldots w_r$. The path P' has smaller number of vertices than path P, and note that the missing root vertex w_1 can be replaced by w_l with the new tooth being w_{l-1} . This gives us a k+2-rake that has smaller number of vertices than the original, which contradicts our minimality assumption.

Next, we show that v_i is not adjacent to w_4, w_5, \ldots, w_r . Again, suppose for contradiction that v_i is adjacent to w_l for some $l = 4, 5, \ldots, r$. Let the path P' be obtained from path P by replacing the subpath $w_1w_2 \ldots w_r$ of P by path $w_1v_iw_lw_{l+1} \ldots w_r$. Again, the path P' has fewer vertices than path P, all the root vertices of P remain in path P', but as v_i is now in the path P', we assign a new tooth w_2 to correspond to the root w_1 . Again, we obtain a k + 2-rake that has smaller number of vertices than the original, a contradiction.

By symmetry, we can show that w_{r+1} is not adjacent to w_l for any $l=2,3,\ldots,r-1$ and v_{i+1} is not adjacent to any of w_1,w_2,\ldots,w_{r-3} . We conclude that none of the wings of the H-graph are adjacent to any of w_4,w_5,\ldots,w_{r-3} . In other words, vertices w_4,w_5,\ldots,w_{r-3} are of degree 2 in G. If w_4w_5 is a cut-edge of G, we have that no vertex of $\{w_0,w_1,w_2,w_3,v_i\}$ is adjacent to any of the vertex of $\{w_{r-2},w_{r-1},w_r,w_{r+1},v_{i+1}\}$. Let $l\leq 3$ be the largest possible such that w_l has degree at least 3 in G, $p\geq r-2$ the smallest possible such that w_p has degree at least 3 in G. Taking the path $w_lw_{l+1}\ldots w_p$ together with another two neighbours of w_l and w_p provides us with an induced H-graph whose base $w_lw_{l+1}\ldots w_p$ has at least s+1 vertices. On the other hand, if w_4w_5 is not a cut-edge in G, then there is a chordless cycle in G containing the edge w_4w_5 and hence this cycle

must contain $w_3w_4w_5...w_{r-2}$ (because of vertices of degree 2). Therefore, we obtain an induced cycle of G with at least $r-4 \ge s+1$ vertices. Hence in both cases we obtain a canonical graph of order at least s as an induced subgraph. This finishes the proof.

3.4. Dense rake subgraphs

Lemma 3. For every s, q and ℓ , there is a number $D = D(s, q, \ell)$ such that every graph containing an ℓ -dense D-rake as a subgraph contains either

- a canonical graph of order at least s as an induced subgraph or
- a biclique of order q as a subgraph.

Proof. To define the number $D = D(s, q, \ell)$, we introduce intermediate notations as follows: $b := 2(q-1)s^q + 2sq + 4$ and $c := R(2, 2, \max(b, 2q))$, where R is the Ramsey number. With these notations the number D is defined as follows: $D = D(s, q, \ell) := Z(\ell c^2, 2q, q)$, where Z is the number defined in Theorem 2.

Consider a graph G containing an ℓ -dense D-rake R^0 as a subgraph. The base of this rake is a path P^0 of length at least D and hence, by Theorem 2, the subgraph of G induced by the base contains either a biclique of order at least q as a subgraph (in which case we are done) or an *induced* path Pof length at least ℓc^2 . Let us call any (inclusionwise) maximal sequence of consecutive vertices of P^0 that belong to P a block. Assume the number of blocks is more than c. Let P' be the subpath of P induced by the first cblocks. Let w_1, \ldots, w_c be the rightmost vertices of the blocks. Let v_1, \ldots, v_c be the vertices such that each v_i is the vertex of P_0 immediately following w_i . Then P' together with v_1, \ldots, v_c create a c-rake with P' being the induced base, v_1, \ldots, v_c being the teeth and w_1, \ldots, w_c being the respective roots. If the number of blocks is at most c, then P^0 must contain a block of size at least ℓc , in which case this block also forms an induced base of a c-rake (since R^0 is ℓ -dense). We see that in either case G has a c-rake with an induced base. According to the definition of c, the c teeth of this rake induce a graph which has either a clique of size 2q (and hence a biclique of order q in which case we are done), or an independent set of size b. By ignoring the teeth outside this set we obtain a b-rake R with an induced base and with teeth forming an independent set.

Let us denote the base of R by U, its vertices by u_1, \ldots, u_m (in the order of their appearances in the path), and the teeth of R by t_1, \ldots, t_b (following the order of their root vertices).

Denote $r := (q-1)s^q + 2$ and consider two sets of teeth $T_1 = \{t_2, t_3, \dots, t_r\}$ and $T_2 = \{t_{b-1}, t_{b-2}, \dots, t_{b-r+1}\}$. By definition of r and b, there are 2sq other teeth between t_r and t_{b-r+1} , and hence there is a set M of 2sq consecutive vertices of U between the root of t_r and the root of t_{b-r+1} . We partition M into 2q subsets (of consecutive vertices of U) of size s each and for $i = 1, \dots, 2q$ denote the i-th subset by M_i .

If each vertex of T_1 has a neighbour in each of the first q sets M_i , then by the Pigeonhole Principle there is a biclique of order q with q vertices in T_1 and q vertices in M. Similarly, a biclique of order q arises if each vertex of T_2 has a neighbour in each of the last q sets M_i . Therefore, we assume that there are two vertices $t_a \in T_1$ and $t_b \in T_2$ and two sets M_x and M_y with x < y such that t_a has no neighbours in M_x , while t_b has no neighbours in M_y .

By definition, t_a has a neighbour in U (its root) on the left of M_x . If additionally t_a has a neighbour to the right of M_x , then a chordless cycle of length at least s arises (since $|M_x| = s$ and t_a has no neighbours in M_x), in which case the lemma is true. This restricts us to the case, when all neighbours of t_a in U are located to the left of M_x . By analogy, we assume that all neighbours of t_b in U are located to the right of M_y . Let u_i be the rightmost neighbour of t_a in U and u_j be the leftmost neighbour of t_b in U. According to the above discussion, i < j and j - j > 2s. But then the vertices $t_a, t_b, u_{i-1}, u_i, \ldots, u_j, u_{j+1}$ induce an H-graph (possibly tight or semi-tight) of order more than s (the existence of vertices u_{i-1} and u_{j+1} follows from the fact that T_1 does not include t_1 , while T_2 does not include t_b).

3.5. Proof of Theorem 1

Combining the results of Lemma 1, Lemma 2 and Lemma 3, we conclude that for every s, q, there is a number X = X(s, q) such that every graph of path-width at least X contains either

- a canonical graph of order at least s as an induced subgraph or
- a biclique of order q as a subgraph.

From this it is not hard to conclude that a class of graphs with unbounded path-width that excludes a biclique of order q must contain an infinite family of distinct canonical graphs, hence the class must be not well-quasi-ordered. Therefore, well-quasi-ordered classes that exclude a biclique of order q for some q, must be of bounded path-width, as required.

References

- [1] J. Balogh, B. Bollobás and D. Weinreich. The speed of hereditary properties of graphs. *Journal of Combinatorial Theory, Series B* 79(2): 131–156, 2000. MR1769217
- [2] B. Courcelle, J. A. Makowsky and U. Rotics. Linear time solvable optimization problems on graphs of bounded clique-width. *Theory Comput. Syst.*, 33(2): 125–150, 2000. MR1739644
- [3] J. Daligault, M. Rao and S. Thomassé. Well-Quasi-Order of Relabel Functions. *Order*, 27: 301–315, 2010. MR2728737
- [4] R. Diestel, Graph Theory. Third edition. Graduate Texts in Mathematics, 173. Springer-Verlag, Berlin, 2005. xvi+411 pp. MR2159259
- [5] G. Ding. On canonical antichains, *Discrete Mathematics*, 309: 1123–1134, 2009. MR2493532
- [6] A. Finkel and Ph. Schnoebelen. Well-structured transition systems everywhere! *Theor. Comput. Sci.*, 256(1–2): 63–92, 2001. MR1821455
- [7] F. Galvin, I. Rival and B. Sands. A Ramsey-Type Theorem for Traceable Graphs. *Journal of Combinatorial Theory, Series B*, 33: 7–16, 1982. MR0678167
- [8] J. B. Kruskal. The theory of well-quasi-ordering: A frequently discovered concept. J. Comb. Theory, Ser. A, 13(3): 297–305, 1972. MR0306057
- [9] V. Lozin, I. Razgon and V. Zamaraev. Well-quasi-ordering versus clique-width. J. Comb. Theory, Ser. B, 130: 1–18, 2018. MR3772732
- [10] S. Norine, P. Seymour, R. Thomas and P. Wollan. Proper minor-closed families are small. J. Comb. Theory Ser. B, 96(5): 754–757, 2006. MR2236510
- [11] T. Kövári, V.T. Sós and P. Turán. On a problem of K. Zarankiewicz. Colloq. Math., 3: 50–57, 1954. MR0065617
- [12] N. Robertson and P.D. Seymour. Graph Minors. XX. Wagner's conjecture, Journal of Combinatorial Theory Ser. B, 92: 325–357, 2004. MR2099147
- [13] N. Robertson and P. Seymour. Graph Minors I. Excluding a Forest. J. Comb. Theory Ser. B, 35: 39–61, 1983. MR0723569

AISTIS ATMINAS
MATHEMATICS INSTITUTE
UNIVERSITY OF WARWICK
COVENTRY CV4 7AL

UK

 $E ext{-}mail\ address: aistis.atminas@gmail.com}$

Vadim Lozin

MATHEMATICS INSTITUTE UNIVERSITY OF WARWICK COVENTRY CV4 7AL

UK

E-mail address: V.Lozin@warwick.ac.uk

IGOR RAZGON

DEPARTMENT OF COMPUTER SCIENCE AND INFORMATION SYSTEMS

BIRKBECK UNIVERSITY OF LONDON

UK

E-mail address: igor@dcs.bbk.ac.uk

RECEIVED 26 JANUARY 2014